ISO/IEC DTR 13211–1:2006
New built-in flags, predicates, and functions
proposal

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Introduction
This proposal specifies a set of built-in predicates and flags to be added to Part 1 of the International Standard for Prolog, ISO/IEC 13211. When evaluating this proposal, please comment each predicate individually by presenting your arguments for either accepting or rejecting its inclusion in the next revision of the Part 1 standard.

This proposal is written as an extension to the ISO/IEC 13211–1 Prolog standard, adopting a similar structure. Specifically, this proposal either adds new sections and clauses to, or modifies the reading of existing clauses on ISO/IEC 13211–1.

This draft proposal may contain in several places informative text, type-set in *italics*. Such informative text is used for editorial comments deemed useful during the development of this draft and may not be included in the final version.

Contributors
This list includes so far the people present at the WG17 ISO meetings collocated with the ICLP’06 and ICLP’09 conferences and people participating on the mailing list discussions.

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1 Scope

This proposal is designed to promote the applicability and portability of Prolog by adding to ISO/IEC 13211–1:1995 a set of built-in predicates and flags that are either common practice and implemented in most Prolog systems or are needed to clarify implementation-dependent behavior. As such, this proposal includes specifications for:

a) A set of flags allowing a programmer to query a system about the floating-point arithmetic implementation and to declare the default encoding for Prolog text

b) Commonly used term testing predicates already available in most Prolog systems that are missing from ISO/IEC 13211–1:1995

c) Commonly used meta-predicates which should be available as built-in predicates in order to provide adequate performance

d) A set of built-in predicates for list processing, providing functionality similar to the atomic processing built-in predicates present on ISO/IEC 13211–1:1995
e) Commonly used evaluable functors that are missing from ISO/IEC 13211–1:1995

NOTE — This part of ISO/IEC 13211 will eventually merge with ISO/IEC 13211–1:1995 resulting in a new version of the Part 1 standard.

3 Definitions

3.34 closure: An atom or a compound term.

6 Syntax

6.4 The operator table

The infix operator ‘><’ (bitwise exclusive or) is added to the operator table with the same specification as the bitwise or and and bitwise and operators:

<table>
<thead>
<tr>
<th>Priority</th>
<th>Specifier</th>
<th>Operators(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>500</td>
<td>yfx</td>
<td>+ - / \ / &gt;&lt;</td>
</tr>
</tbody>
</table>

The prefix operator ‘+’ (unary plus) is added to the operator table with the same specification as the unary minus or the bitwise complement operators:

<table>
<thead>
<tr>
<th>Priority</th>
<th>Specifier</th>
<th>Operators(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>200</td>
<td>fy</td>
<td>+ - \</td>
</tr>
</tbody>
</table>

The infix operator ‘^’ is used in arithmetic expressions to represent integer exponentiation.

7 Language concepts and semantics

7.1 Types

7.1.6 Related terms

7.9 Pair

P is a pair if it is a compound term ‘-’(Key, Value) where Key and Value are terms.

NOTE — In Prolog text and this part of ISO/IEC 13211 a pair ‘-’(Key, Value) is normally written as Key=Value or (Key)=Value depending on whether or not Key and Value are operators.
7.4 Prolog text
7.4.2 Directives

7.10 encoding/1

A directive \texttt{encoding(E)} specifies that the Prolog text being prepared for execution uses the encoding \texttt{E}. When used, this directive shall be the first term, on the first line, in a Prolog text with no extra layout characters (6.5.4) before the \texttt{:\=/1} directive operator. Moreover, a single layout character shall be used between the directive operator and the directive functor. No layout characters or comments shall appear between the directive opening and closing parenthesis.

7.11 if/1, elif/1, else/0, endif/0

The directives \texttt{if(Goal)}, \texttt{elif(Goal)}, \texttt{else}, and \texttt{endif} support the conditional of source code.

7.8 Control constructs

7.8.11 call/2-N

7.1 Description

call(Closure, Arg1, ...) is true iff call(Goal) is true where \texttt{Goal} is constructed by appending \texttt{Arg1}, ... additional arguments to the arguments (if any) of \texttt{Closure}.

7.2 Template and modes

call(@callable\_term, \texttt{?term}, ...)

7.3 Errors

a) \texttt{Closure} is a variable
   \begin{itemize}
   \item \texttt{instantiation\_error}
   \end{itemize}

b) \texttt{Closure} is neither a variable nor a callable term
   \begin{itemize}
   \item \texttt{type\_error(callable, Closure)}
   \end{itemize}

c) The number of arguments in the resulting goal exceeds the implementation defined maximum arity for compound terms
   \begin{itemize}
   \item \texttt{representation\_error(max\_arity)}
   \end{itemize}

7.4 Examples
### Table 1: Status of current Prolog implementations providing built-in support for the call/2-N control construct.

<table>
<thead>
<tr>
<th>System</th>
<th>N</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>B-Prolog</td>
<td>10/65535</td>
<td>(interpreter/compiler i.e. maximum arity)</td>
</tr>
<tr>
<td>CxProlog</td>
<td>9</td>
<td></td>
</tr>
<tr>
<td>GNU Prolog</td>
<td>11</td>
<td></td>
</tr>
<tr>
<td>SICStus Prolog</td>
<td>255</td>
<td>(maximum arity)</td>
</tr>
<tr>
<td>SWI-Prolog</td>
<td>8</td>
<td>(meta.predicate/1 directive limited to predicates of arity up to 8)</td>
</tr>
<tr>
<td>XSB</td>
<td>11</td>
<td></td>
</tr>
<tr>
<td>YAP</td>
<td>12</td>
<td></td>
</tr>
</tbody>
</table>

```prolog
call(integer, 3).
   succeeds.

call(atom_concat(pro), log, Atom).
   succeeds, unifying Atom with prolog.
```

#### 7.8.12 setup_call_cleanup/3

For details on the specification proposal for this built-in predicate please check: [http://www.complang.tuwien.ac.at/ulrich/iso-prolog/cleanup](http://www.complang.tuwien.ac.at/ulrich/iso-prolog/cleanup)

#### 7.10 Input/output

##### 7.10.3 Standard Streams

*Added user_error alias.*

Three streams are predefined and open during the execution of every goal: the standard input stream has the alias `user_input`, the standard output stream has the alias `user_output`, and the standard error output stream has the alias `user_error`.

**NOTES**

3 Prolog implementation on systems that do not support an error output stream shall redirect output to the standard output stream.

#### 7.11 Options on stream creation

An implementation may optionally support the following stream-options:
bom(Bool) — If Bool (7.1.4.2) is true then a Unicode encoding Byte Order Mark shall be written when opening a text stream for writing in mode write or is probed for when opening the text stream for reading. This option shall be ignored when opening a stream in mode append. If Bool is false then a Unicode encoding Byte Order Mark shall not be written when opening the text stream for writing.

When no bom(Bool) stream-option is specified, the default value shall be true when the text stream is opened for reading and false when the text stream is opened for writing.

encoding(Encoding) — Encoding is an atom representing the text encoding that shall be used when opening the stream for writing or the text encoding of the stream opened for reading.

When opening the text stream for reading with the default bom(Bool) stream-option value or by explicitly specifying the bom(true) stream-option, if a Byte Order Mark is detected, it will be used to set the corresponding Unicode text stream encoding, overriding any encoding(Encoding) that might be also specified.

NOTES

1 These stream-options imply the stream-option type(text).

2 The set of supported text encodings is implementation-defined.

7.13 Stream properties

An implementation may optionally support the following stream properties:

bom(Bool) — If present and if Bool (7.1.4.2) is true, a Unicode encoding Byte Order Mark was detected while opening the text stream for reading or a Byte Order Mark was written while opening the text stream for writing.

encoding(Encoding) — Encoding used for the text stream.

NOTE — These stream properties imply the stream property type(text).

7.11 Flags

7.11.2 Flags defining float type F

7.1 Flag: float_mantissa_digits
Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is the number of significant digits on the mantissa of a normalized floating point number (in base 10), an implementation defined integer value.

7.2 Flag: float epsilon

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is the distance from 1.0 to the next largest floating point number, an implementation defined value. Thus, it allows the programmer to query an implementation about the relative accuracy when performing arithmetic with floating point numbers.

7.3 Flag: float min exponent

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is smallest possible value of the exponent of a normalized floating point number, an implementation defined integer value.

7.4 Flag: float max exponent

Possible values: the default value only

Default: implementation defined

Changeable: No

Description: The value of this flag is greatest value of the exponent of a normalized floating point number, an implementation defined integer value.
7.11.3 Other flags

7.6 Flag: unification_subject_to_occurs_check

Possible values: fail, cyclic, unsafe

Default: implementation defined

Changeable: implementation defined

Description: This read-only flag describes the behaviour of the Prolog system when a variable is unified with a compound term that contains it (STO unification, 3.165). The flag value fail implies that STO unification simply fails. The flag value unsafe means that if an STO unification is encountered the further behavior of the system is undefined. The flag value cyclic implies that STO unifications will be successful and result in the creation of cyclic terms. Moreover, this flag indicates that unification between cyclic terms terminates.

NOTE

The flag value fail means that the built-in predicate =/2 (8.2.1) behaves exactly as the built-in predicate unify_with_occurs_check/2 (8.2.2). The flag value unsafe may imply that the STO unification itself, or further unifications or built-in predicate calls may not terminate, or cause the system to fail or raise an exception.

Examples:

| ?- X = f(X). |
| fail: fails |
| cyclic: succeeds and unifies X with a cyclic term f(f(f(...))). |
| unsafe: undefined. Often succeeds, but subsequent use of X, as e.g. in X=X, causes an error. |

| ?- X = f(X), Y = f(Y), X = Y. |
| ?- g(X,Y,X) = g(f(X),f(Y),Y). |
| ?- X = f(X), Y = f(Y), X == Y. |
| ?- X = f(X), asserta(p(X)). |

For all the above four examples:

- fail: fails
- cyclic: succeeds and unifies both X and Y with a cyclic term f(f(f(...))).
7.7 Flag: encoding

Possible values: an implementation defined atom

Default: implementation defined

Changeable: implementation defined

Description: This flag represents the default encoding for text streams. An implementation shall document if the flag value can be changed by programmer as well all the supported encodings. The text encoding names to be used shall be the ones specified by the Internet Assigned Numbers Authority (IANA) and marked as “(preferred MIME name)”:

http://www.iana.org/assignments/character-sets

NOTE — Implementations can define aliases for the standard text encodings names.

7.12 Errors

7.12.2 Error classification

The following types are added to the classification of 7.12.2 of ISO/IEC 13211-1.

a) The list of valid types is extended by the addition of pair (see 7.12.2 b of ISO/IEC 13211-1).

b) The list of valid domains is extended by the addition of order and predicate.property (see 7.12.2 c of ISO/IEC 13211-1).

7.13 Predicate properties

The properties of procedures can be found using the built-in predicate predicate.property(Callable, Property), where Callable is a callable term. The predicate properties supported shall include:

- static — The predicate is static
- dynamic — The predicate is dynamic
- built_in — The predicate is a built-in predicate
- multifile — The predicate is the subject of a multifile directive

A processor may support one or more additional predicate properties as an implementation specific feature. Implementation-defined properties are not required to be atomic terms.
8 Built-in predicates

The following sections extend, with the specified number, the corresponding ISO/IEC 13211–1 sections:

8.2 Term unification

8.2.4 subsumes/2

8.1 Description

\texttt{subsumes(General, Specific)} is true iff there is a substitution \( \theta \), including the empty substitution, such that the term \( \text{General} \) is instantiated to \( \text{General}_\theta = \text{Specific}_\theta \). This predicate provides a one-way unification.

8.2 Template and modes

\texttt{subsumes(?term, @term)}

8.3 Errors

None.

8.4 Examples

\begin{verbatim}
subsumes(f(X,Y), f(Z,Z)).
Succeeds, unifying both X and Y to Z.

subsumes(f(Z,Z), f(X,Y)).
Fails.
\end{verbatim}

8.3 Type testing

8.3.9 callable/1

8.1 Description

\texttt{callable(Term)} is true iff \( \text{Term} \) is a callable term.

8.2 Template and modes

\texttt{callable(@term)}

8.3 Errors

None.
8.4 Examples

callable(a).
   Succeeds.

callable(3).
   Fails.

callable(X).
   Fails.

8.3.10 ground/1

8.1 Description

ground(Term) is true iff Term is a ground term.

8.2 Template and modes

ground(@term)

8.3 Errors

None.

8.4 Examples

   ground(3).
      Succeeds.

   ground(a(1, _)).
      Fails.

8.3.11 acyclic_term/1

8.1 Description

acyclic_term(Term) is true iff Term is an acyclic term. For implementations not supporting STO unification \[7.0\] calls to this predicate simply succeed.

8.2 Template and modes

acyclic_term(@term)
8.3 Errors

None.

8.4 Examples

acyclic_term(a(1, _)).
   Succeeds.

X = f(X), acyclic_term(X).
   Fails.

8.3.12 cyclic_term/1

8.1 Description

cyclic_term(Term) is true iff Term is a cyclic term. For implementations not supporting STO unification \[7.6\], calls to this predicate simply fail.

8.2 Template and modes

cyclic_term(@term)

8.3 Errors

None.

8.4 Examples

   cyclic_term(a(1, _)).
      Fails.

   X = f(X), cyclic_term(X).
      Succeeds.

8.4 Term comparison

8.4.2 compare/3

8.1 Description

compare(Order, Term1, Term2) is true iff Order corresponds to the standard order between Term1 and Term2. The argument Order is unified with the atom < when Term1 is less than Term2, with the atom = when Term1 and Term2 are equal, and with the atom > when Term1 is greater than Term2.
8.2 Template and modes

compare(?atom, @term, @term)

8.3 Errors

a) Order is neither a variable nor an atom
   — type_error(atom, Order)

b) Order an atom other than <, =, or >
   — domain_error(order, Order)

8.4 Examples

compare(Order, 3, 5).
   Succeeds, unifying Order with <.

compare(Order, d, d).
   Succeeds, unifying Order with =.

compare(Order, 3, 3.0).
   Succeeds, unifying Order with >.

8.5 Term creation and decomposition

8.5.5 numbervars/3

8.1 Description

numbervars(Term, Start, End) is true. This predicate unifies each free variable on Term with a compound term with the format ‘$VAR'(N) where N is an integer starting from Start and ending at End−1.

8.2 Template and modes

numbervars(?nonvar, +integer, -integer)

8.3 Errors

a) Start is a variable
   — instantiation_error

b) Start is neither a variable nor an integer
   — type_error(integer, Start)
8.4 Examples

\texttt{numbervars(foo(A, B, A), 0, End).}
Succeeds, unifying A with \texttt{\$VAR'(0)}, B with \texttt{\$VAR'(1)},
and End with 2.

8.8 Clause retrieval and information

8.8.3 \texttt{predicate_property/2}

8.1 Description

\texttt{predicate_property(Head, Property)} is true iff the procedure associated with
the argument \texttt{Head} (3.84) has predicate property \texttt{Property}.

Procedurally, \texttt{predicate_property(Head, Property)} is executed as follows

a) Determines the principal functor \(P\) and arity \(N\) associated with \texttt{Head}. \(P/N\)
is the associated predicate indicator

b) Searches the complete database and creates a set \texttt{SetPP} of all terms \texttt{PP}
such that \(P/N\) identifies a procedure which has predicate property \texttt{PP} and
\texttt{PP} is unifiable with \texttt{Property}

c) If \texttt{SetPP} is non empty set proceeds to 8.1 e,

d) Else the goal fails

e) Chooses the first element \texttt{PPP} of \texttt{SetPP}, unifies \texttt{PPP} with \texttt{Property} and
the predicate succeeds

f) If all elements of \texttt{SetPP} have been chosen the predicate fails

g) Else chooses the first element \texttt{PPP} of \texttt{SetPP} that has not already been
chosen, unifies \texttt{PPP} with \texttt{Property} and the goal succeeds

\texttt{predicate_property(Head, Property)} is re-executable. On backtracking, continue at 8.1 f.

The order in which properties are found by \texttt{predicate_property/2} is implement-
ation dependent.

NOTES

1 \hspace{1cm} A processor may support, as an implementation specific feature, additional
predicate properties.

2 \hspace{1cm} For a dynamic predicate, all proprieties related to its definition shall be
removed when the predicate is abolished.
8.2 Template and modes

predicate_property(?callable_term, ?predicate_property)

8.3 Errors

a) Head is neither a variable nor a callable term
   — type_error(callable, Head)

b) Property is neither a variable nor a predicate property
   — domain_error(predicate_property, Property)

8.4 Examples

predicate_property(once(_), built_in).
   Succeeds.

predicate_property(atom_codes(_, _), Property).
   Succeeds unifying Property with static.
   On re-execution, succeeds unifying Property with built_in.

8.9 Clause creation and destruction

8.9.3 retract/1

8.1 Errors

There is a typo on the current standard in the specification of the permission_error
exception that should use the atom modify instead of access in order to match
the specification of other database predicates.

This typo may or may not be already corrected in the 13211-1 errata recently
published.

a) ...

b) ...

c) The predicate indicator Pred of Head is that of a static procedure
   — permission_error(modify, static_procedure, Pred)

8.9.5 retractall/1

8.1 Description

retractall(Head) is true. Procedurally, the predicate shall behave as if defined by:
retractall(Head) :-
    retract((Head :- _)),
    fail.
retractall(_).

Procedurally, retractall(Head) is executed as follows:

a) Determines the principal functor P and arity N associated with Head. P/N is the associated predicate indicator
b) If the database contains a dynamic procedure whose predicate indicator is P/N, then proceeds to 8.1 d,
c) Else the goal succeeds.
d) Retracts from the database all clauses whose head unifies with Head and the goal succeeds

NOTES

1 The dynamic predicate shall continue to be known by the system even when all of its clauses are removed.

2 This predicate does not change any of the standard predicate properties of the referenced predicate (as reported by predicate_property(Head, Property)), even when all of its clauses are removed.

8.2 Template and modes

retractall(@callable_term)

8.3 Errors

a) Head is a variable
   — instantiation_error
b) Head is neither a variable nor a callable term
   — type_error(callable, Generate)
c) The predicate indicator Pred of Head is that of a static procedure
   — permission_error(modify, static_procedure, Pred)

8.4 Examples

The examples defined in this subclause assume the database has been created from the following Prolog text:
:- dynamic(insect/1).
insect(ant).
insect(bee).
insect(spider).

retractall(insect(bee)).
    Succeeds, retracting the clause 'insect(bee)'.

retractall(insect(_)).
    Succeeds, retracting all the clauses of predicate insect/1.

retractall(insect(elephant)).
    Succeeds.

retractall(mammal(_)).
    Succeeds.

retractall(3).
    type_error(callable, 3)

8.10 All solutions

8.10.4 forall/2

8.10.4.1 Description

forall(Generate, Test) is true iff for all possible bindings of Generate, the goal Test is true. Procedurally, the predicate shall behave as if defined by \+ (call(Generator), \+ call(Test)).

8.10.4.2 Template and modes

forall(@callable_term, @callable_term)

8.10.4.3 Errors

a) Generate is a variable
   — instantiation_error

b) Generate is neither a variable nor a callable term
   — type_error(callable, Generate)

c) Test is a variable
   — instantiation_error

d) Test is neither a variable nor a callable term
   — type_error(callable, Test)
8.10.4.4 Examples

The following examples assume that the predicate a/1 and b/1 are defined with the following clauses:

```
a(1). a(2). a(3).
b(1, a). b(2, b). b(3, c).
```

forall(fail, true).
Succeeds.

forall(a(X), b(X, _)).
Succeeds.

forall(a(X), b(_, X)).
Fails.

forall(b(_, Y), write(Y))
Succeeds, outputting the characters
abc
to the current output stream.

8.10.5 sort/2

8.10.5.1 Description

sort(List, Sorted) is true iff Sorted is a list containing the non-duplicated elements of List sorted in ascending order following standard order (7.2).

8.10.5.2 Template and modes

sort(@list, ?list)

8.10.5.3 Errors

a) List is a partial list
   — instantiation_error

b) List is neither a partial list nor a list
   — type_error(list, List)

c) Sorted is neither a partial list nor a list
   — type_error(list, Sorted)

8.10.5.4 Examples

```
sort([1, 2, 1, 8, 4], Sorted).
Succeeds, unifying Sorted with [1, 2, 4, 8].
```
8.10.6 keysort/2

8.10.6.1 Description

`keysort(List, Sorted)` is true iff `List` is a list of elements with the format `Key-Value` and `Sorted` is a list containing the elements of `List` sorted according to the value of `Key` in ascending order following standard order (7.2). The sorting procedure shall be stable. I.e. the relative order of elements of `List` with the same key shall not change in the `Sorted` list.

8.10.6.2 Template and modes

`keysort(@list, ?list)`

8.10.6.3 Errors

a) `List` is a partial list
   — `instantiation_error`

b) `List` is neither a partial list nor a list
   — `type_error(list, List)`

c) An element `Element` of `List` is a variable
   — `instantiation_error`

d) An element `Element` of `List` is neither a variable nor a `'-'/2` compound term
   — `type_error(pair, Element)`

e) `Sorted` is neither a partial list nor a list
   — `type_error(list, Sorted)`

8.10.6.4 Examples

`keysort([1-a, 3-f(_,), 1-z, 2-44], Sorted).`
`Succeeds unifying Sorted with [1-a, 1-z, 2-44, 3-f(_)].`

9 Evaluable functors

9.1 The simple arithmetic functors

*The unary plus evaluable arithmetic functor is added.*

9.1.1 Evaluable functors and operations

<table>
<thead>
<tr>
<th>Evaluable functor</th>
<th>Operation</th>
</tr>
</thead>
<tbody>
<tr>
<td>(+)/1</td>
<td>pos_i, pos_f</td>
</tr>
</tbody>
</table>
9.1.3 Integer operations and axioms

The following operations are specified:

\[ pos_I : I \rightarrow I \]

For all \( x \in I \), the following axioms shall apply:

\[ pos_I(x) = x \]

9.1.4 Floating point operations and axioms

The following operations are specified:

\[ pos_F : F \rightarrow F \]

For all \( x \in F \), the following axioms shall apply:

\[ pos_F(x) = x \]

9.3 Other arithmetic functors

9.3.8 \(^*/2\)

9.3.8.1 Description

\(^*(VI, VJ)\) evaluates the expression \( I \) with value \( VI \), the expression \( J \) with value \( VJ \), and has the value of the integer exponentiation of \( VI \) to \( VJ \).

9.3.8.2 Template and modes

\(^*(\text{int-exp}, \text{int-exp}) = \text{integer} \)

9.3.8.3 Errors

a) \( I \) is a variable
   — \text{instantiation_error}

b) \( I \) is not a variable and \( VI \) is not an integer
   — \text{type_error(integer, VI)}

c) \( J \) is a variable
   — \text{instantiation_error}

d) \( J \) is not a variable and \( VJ \) is not an integer
   — \text{type_error(integer, VJ)}
9.3.8.4 Examples

3412.
Evaluates to 2386420683693101056.

9.3.9 log/2

9.3.9.1 Description

log(B, X) evaluates the expression B with value \( V_B \), the expression X with value \( V_X \), and has the value of the logarithm to base \( V_B \) of \( V_X \).

9.3.9.2 Template and modes

\[
\begin{align*}
\text{log(int-exp, float-exp)} &= \text{float} \\
\text{log(int-exp, int-exp)} &= \text{float}
\end{align*}
\]

9.3.9.3 Errors

a) \( B \) is a variable
   — instantiation_error

b) \( B \) is not a variable and \( V_B \) is not an integer
   — type_error(integer, \( V_B \))

c) \( V_B \) is zero or negative
   — evaluation_error(undefined)

d) \( X \) is a variable
   — instantiation_error

e) \( X \) is not a variable and \( V_X \) is not a number
   — type_error(number, \( V_X \))

f) \( V_X \) is zero or negative
   — evaluation_error(undefined)

9.3.9.4 Examples

\[
\begin{align*}
\text{log(10, 10.0)}.
\text{Evaluates to 1.0.}
\end{align*}
\]

9.3.10 gcd/2

9.3.10.1 Description

\( \text{gcd}(I, J) \) evaluates the expression I with value \( V_I \), the expression J with value \( V_J \), and has the value of the greatest common divisor of \( V_I \) of \( V_J \).
9.3.10.2 Template and modes

\[ \text{gcd}(\text{int-exp}, \text{int-exp}) = \text{integer} \]

9.3.10.3 Errors

a) \( I \) is a variable  
   \[ \text{instantiation\_error} \]

b) \( I \) is not a variable and \( VI \) is not an integer  
   \[ \text{type\_error(integer, VI)} \]

c) \( J \) is a variable  
   \[ \text{instantiation\_error} \]

d) \( J \) is not a variable and \( VJ \) is not an integer  
   \[ \text{type\_error(integer, VJ)} \]

9.3.10.4 Examples

\[ \text{gcd}(2, 3). \]
   Evaluates to 1.

9.3.11 \( \text{max/2} \)

9.3.11.1 Description

\( \text{max}(X, Y) \) evaluates the expression \( X \) with value \( VX \), the expression \( Y \) with value \( VY \), and has the value of the maximum of \( VX \) and \( VY \). When used with expressions of mixed-types, the result is implementation-dependent; an implementation may chose either to return a value or to throw an exception.

9.3.11.2 Template and modes

\[ \text{max(float-exp, float-exp)} = \text{float} \]
\[ \text{max(float-exp, int-exp)} = \text{implementation-dependent result} \]
\[ \text{max(int-exp, float-exp)} = \text{implementation-dependent result} \]
\[ \text{max(int-exp, int-exp)} = \text{integer} \]

9.3.11.3 Errors

a) \( X \) is a variable or \( Y \) is a variable  
   \[ \text{instantiation\_error} \]

b) \( X \) is not a variable and \( VX \) is not a number  
   \[ \text{type\_error(number, VX)} \]

c) \( Y \) is not a variable and \( VY \) is not a number  
   \[ \text{type\_error(number, VY)} \]
9.3.11.4 Examples

max(2, 3)
   Evaluates to 3.

max(2.0, 3.0)
   Evaluates to 3.0.

max(0, 0.0).
   Implementation-dependent result.

9.3.12 min/2

9.3.12.1 Description

\[\text{min}(X, Y)\] evaluates the expression \(X\) with value \(V_X\), the expression \(Y\) with value \(V_Y\), and has the value of the minimum of \(V_X\) and \(V_Y\). When used with expressions of mixed-types, the result is implementation-dependent; an implementation may chose either to return a value or to throw an exception.

9.3.12.2 Template and modes

\[
\begin{align*}
\text{min(float-exp, float-exp)} &= \text{float} \\
\text{min(float-exp, int-exp)} &= \text{implementation-dependent result} \\
\text{min(int-exp, float-exp)} &= \text{implementation-dependent result} \\
\text{min(int-exp, int-exp)} &= \text{integer}
\end{align*}
\]

9.3.12.3 Errors

a) \(X\) is a variable or \(Y\) is a variable
   — instantiation_error

b) \(X\) is not a variable and \(V_X\) is not a number
   — type_error(number, \(V_X\))

c) \(Y\) is not a variable and \(V_Y\) is not a number
   — type_error(number, \(V_Y\))

9.3.12.4 Examples

\[\text{min}(2, 3)\]
   Evaluates to 2.

\[\text{min}(2.0, 3.0)\]
   Evaluates to 2.0.

\[\text{min}(0, 0.0)\]
   Implementation-dependent result.
9 EVALUABLE FUNCTORS

9.4 Bitwise functors

9.4.6 (\texttt{\textgreater\textless})/2 – bitwise exclusive or

9.4.6.1 Description

\texttt{\textgreater\textless}'(B1, B2) evaluates the expressions \(B1\) and \(B2\) with values \(VB1\) and \(VB2\) and has the value such that each bit is set iff only one of the corresponding bits in \(VB1\) and \(VB2\) is set.

The value shall be implementation defined if \(VB1\) or \(VB2\) is negative.

9.4.6.2 Template and modes

\texttt{\textgreater\textless}'(\texttt{int-exp, int-exp}) = integer

NOTE — \texttt{\textgreater\textless} is an infix predefined operator (see 6.3.4.4).

9.4.6.3 Errors

a) \(B1\) is a variable
   — instantiation_error

b) \(B2\) is a variable
   — instantiation_error

c) \(B1\) is not a variable and \(VB1\) is not an integer
   — type_error(integer, VB1)

d) \(B2\) is not a variable and \(VB2\) is not an integer
   — type_error(integer, VB2)

9.4.6.4 Examples

\texttt{\textgreater\textless}'(10, 12).
  Evaluates to the value 6.

\texttt{\textgreater\textless}'(125, 255).
  Evaluates to to the value 130.

\texttt{\textgreater\textless}'(-10, 12).
  Evaluates to an implementation defined value.

\texttt{\textgreater\textless}'(77, \texttt{N})
  instantiation_error.

\texttt{\textgreater\textless}'(\texttt{foo}, 2)
  type_error(integer, foo).
9.5 Trigonometric functors

Assuming that we will be adding the trigonometric functions described below to the revised core standard, it’s probably best if we gather all the trigonometric functions under their own section. The following order is assumed below: sin/1, cos/1, tan/1, asin/1, acos/1, atan/1, and atan/2.

9.5.3 tan/1

9.5.3.1 Description

tan(X) evaluates the expression X with value VX and has the value of the tangent of VX (measured in radians).

9.5.3.2 Template and modes

\[
\text{tan(float-exp)} = \text{float} \\
\text{tan(int-exp)} = \text{float}
\]

9.5.3.3 Errors

a) X is a variable
   — instantiation_error

b) X is not a variable and VX is not a number
   — type_error(number, VX)

9.5.3.4 Examples

\[
\text{tan(pi)}.
\]

Evaluates to 0.0.

9.5.4 asin/1

9.5.4.1 Description

asin(X) evaluates the expression X with value VX and has the value of the arc sine of VX (in radians).

9.5.4.2 Template and modes

\[
\text{asin(float-exp)} = \text{float} \\
\text{asin(int-exp)} = \text{float}
\]
9.5.4.3 Errors

a) $X$ is a variable
   — instantiation_error

b) $X$ is not a variable and $VX$ is not a number
   — type_error(number, VX)

9.5.4.4 Examples

asin(1.0).
   Evaluates to a value approximately equal to 1.570796326795.

9.5.5 acos/1

9.5.5.1 Description

$acos(X)$ evaluates the expression $X$ with value $VX$ and has the value of the arc
  cosine of $VX$ (in radians).

9.5.5.2 Template and modes

acos(float-exp) = float
acos(int-exp) = float

9.5.5.3 Errors

a) $X$ is a variable
   — instantiation_error

b) $X$ is not a variable and $VX$ is not a number
   — type_error(number, VX)

9.5.5.4 Examples

acos(0.0).
   Evaluates to a value approximately equal to 1.570796326795.

9.5.6 atan/2

9.5.6.1 Description

$atan(Y, X)$ evaluates the expression $Y$ with value $VY$, the expression $X$ with
  value $VX$, and computes the principal value of the arc tangent of $VY/VX$ (in
  radians), using the signs of both arguments to determine the quadrant of the
  return value. When both arguments are 0.0, an implementation-dependent
  floating-point value is returned.
9.5.6.2 Template and modes

atan(float-exp, float-exp) = float
atan(float-exp, int-exp) = float
atan(int-exp, float-exp) = float
atan(int-exp, int-exp) = float

9.5.6.3 Errors

a) X is a variable
   — instantiation_error

b) Y is not a variable and VY is not a number
   — type_error(number, VY)

c) X is not a variable and VX is not a number
   — type_error(number, VX)

9.5.6.4 Examples

atan(0.0, -0.0).
   Evaluates to a value approximately equal to 3.14159265358979.

9.6 Mathematical constants

9.6.1 pi/0

9.6.1.1 Description

pi evaluates to the floating-point number which best approximates the mathematical constant π, the ratio of a circle’s circumference to its diameter.

9.6.1.2 Examples

pi.
   Evaluates to the corresponding mathematical constant.
   The accuracy of the result is implementation-defined.

9.6.2 e/0

9.6.2.1 Description

e evaluates to the floating-point number which best approximates the mathematical constant e, the base of natural logarithms.

9.6.2.2 Examples

e.
   Evaluates to the corresponding mathematical constant.
   The accuracy of the result is implementation-defined.
9.6.3 epsilon/0

9.6.3.1 Description

epsilon evaluates to the distance from 1.0 to the next largest floating point number, an implementation defined value. Thus, it allows the programmer to retrieve the relative accuracy when performing arithmetic with floating point numbers.

9.6.3.2 Examples

epsilon.

Evaluates to an implementation defined value.